CLAIMS:

1. A method of generating a linear transformation matrix A for use in a symmetric-key cipher, the method including:

- generating a binary [n,k,d] error-correcting code, represented by a generator matrix $G \in \mathbf{Z}_2^{k\times n}$ in a standard form $G = (I_k \,||\, B)$, with $B \in \mathbf{Z}_2^{k\times(n-k)}$, where $k < n < 2k$, and d is the minimum distance of the binary error-correcting code;
- extending matrix B with 2k-n columns such that a resulting matrix C is non-singular, and
- deriving matrix A from matrix C.

2. A method as claimed in claim 1, wherein the step of extending matrix B with 2k-n columns includes:
in an iterative manner:

- (pseudo-)randomly generating 2k-n columns, each with k binary elements;
- forming a test matrix consisting of the n-k columns of B and the 2k-n generated columns; and
- checking whether the test matrix is non-singular,

until a non-singular test matrix has been found; and
using the found test matrix as matrix C.

3. A method as claimed in claim 1, wherein the step of deriving matrix A from matrix C includes:

- determining two permutation matrices $P_1,P_2 \in \mathbf{Z}_2^{k\times k}$ such that all codewords in an [2k,k,d] error-correcting code, represented by the generator matrix $(I \,||\, P_1\, C\, P_2)$, have a predetermined multi-bit weight; and
- using $P_1\, C\, P_2$ as matrix A.

4. A method as claimed in claim 3, wherein the cipher includes a round function with an S-box layer with S-boxes operating on m-bit sub-blocks, and the minimum

predetermined multi-bit weight over all non-zero codewords equals a predetermined m-bit weight.

5. A method as claimed in claim 3, wherein the step of determining the two permutation matrices P_1 and P_2 includes iteratively generating the matrices in a (pseudo-) random manner.

6 A method as claimed in claim 1, wherein the cipher includes a round function operating on 32-bit blocks and wherein the step of generating a [n,k,d] error-correcting code includes:

generating a binary extended Bose-Chaudhuri-Hocquenghem (XBCH) [64, 36, 12] code; and

shortening this code to a [60, 32, 12] shortened XBCH code by deleting four rows.

7. A computer program product, wherein the program product is operative to cause a processor to perform the method of claim 1.

8. A system for cryptographically converting an input data block into an output data block; the data blocks comprising n data bits; the system including:

- an input for receiving the input data block;
- a storage for storing a linear transformation matrix A, generated according to the method of claim 1,
- a cryptographic processor performing a linear transformation on the input data block or a derivative of the input data block using the linear transformation matrix A; and
- an output for outputting the processed input data block..